



اسم المادة : تقنيات رقمية اسم التدريسي :م.م علياء محد جواد

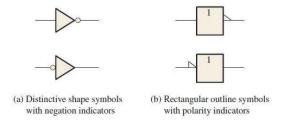
المرحلة : الثانية السنة الدراسية :2024\_2023

عنوان المحاظرة:Logic Gates

1

#### **Logic Gates**

- The inverter (NOT circuit) performs the operation called *inversion* or *complementation*.
- The inverter changes one logic level to the opposite level. In terms of bits, it changes a 1 to a 0 and a 0 to a 1.
- Standard logic symbols for the **inverter** are shown below:



• The "bubble" indicates negation (**inversion** or *complementation*) when it appears on the input or output of any logic element.



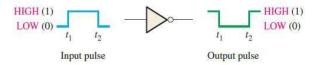


## Logic Gates The inverter (NOT circuit)

Inverter Truth Table

Inverter truth table.				
Input Outpu				
LOW (0)	HIGH (1)			
HIGH (1)	LOW (0)			

- A truth table shows the output for each possible input in terms of levels and corresponding bits.
- Inverter Operation



3

# Logic Gates The inverter (NOT circuit)

- Logic Expression for an Inverter
- In **Boolean algebra**, which is the mathematics of logic circuits, a variable is generally designated by one or two letters although there can be more.
- The **complement** of a variable is designated by a bar over the letter.
- The operation of an inverter (NOT circuit) can be expressed as follows: If the input variable is called A and the output variable is called X, then

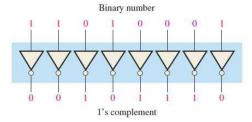
$$X = \overline{A}$$
  $A \longrightarrow X = \overline{A}$ 





# The inverter (NOT circuit) An Application

• The figure below shows a circuit for producing the 1's complement of an 8-bit binary number.

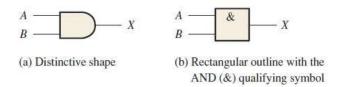


• The bits of the binary number are applied to the inverter inputs and the 1's complement of the number appears on the outputs.

5

#### The AND Gate

- The term gate is used to describe a circuit that performs a basic logic operation.
- The AND gate is one of the basic gates that can be combined to form any logic function.
- An AND gate can have two or more inputs and performs what is known as logical multiplication.



Inp	outs	Output
A	$\boldsymbol{B}$	X
0	0	0
0	1	0
1	0	0
1	1	1





## **Logic Gates**

• The total number of possible combinations of **binary inputs** to a gate is determined by the following formula:

$$N = 2^{n}$$

where N is the number of possible input combinations and n is the number of input variables.

Example: Develop the truth table for a 3-input AND gate

Solution:

	Inputs	Output		
$\boldsymbol{A}$	$\boldsymbol{B}$	C	X	
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	0	
1	0	0	0	
1	0	1	0	
1	1	0	0	
1	1	1	1	

7

#### **AND Gate**

- Logic Expressions for an AND Gate
- The logical AND function of two variables is represented mathematically either by placing a dot between the two variables, as A.B, or by simply writing the adjacent letters without the dot, as AB.
- Boolean multiplication follows the same basic rules governing binary multiplication, and are as follows:

 $0 \cdot 0 = 0$  $0 \cdot 1 = 0$  $1 \cdot 0 = 0$  $1 \cdot 1 = 1$ 

Boolean multiplication is the same as the AND function.



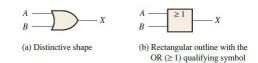
Boolean expressions for AND gates with two, three, and four inputs.





#### The OR Gate

- An OR gate performs what is known as logical addition.
- An OR gate can have two or more inputs and one output.



• OR Gate Truth Table
Truth table for a 2-input

OR g	ate.	W 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
Inp	outs	Output
$\boldsymbol{A}$	$\boldsymbol{B}$	X
0	0	0
0	1	1
1	0	1
1	1	1

С

#### **OR Gate**

- Logic Expressions for an OR Gate
- The logical OR function of two variables is represented mathematically by a + between the two variables, e.g., X= A + B. The plus sign is read as "OR."
- The basic rules for **Boolean addition** are as follows:

$$0 + 0 = 0$$
  
 $0 + 1 = 1$   
 $1 + 0 = 1$   
 $1 + 1 = 1$ 

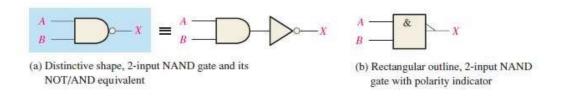
Boolean addition is the same as the OR function.





#### **NAND** Gate

- The NAND gate is a popular logic element because it can be used as a universal gate such that NAND gates can be used in combination to perform the AND, OR, and inverter operations.
- The NAND gate is the same as the AND gate except the output is
- inverted.



11

#### **NAND** Gate

- A **NAND** gate produces a LOW output only when all the inputs are HIGH. When any of the inputs is LOW, the output will be HIGH.
- The table below is for the specific case of a 2-input NAND gate

NAND gate.							
Inp	outs	Output					
A	В	X					
0	0	1					
0	1	1					
1	0	1					

Truth table for a 2-input

1 = HIGH, 0 = LOW





#### **NAND** Gate

- Negative-OR Equivalent Operation of a NAND Gate
- For a 2-input NAND gate performing a negative-OR operation, output X is HIGH when either input A or input B is LOW, or when both A and B are LOW.



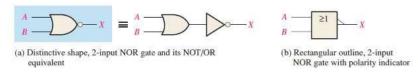
- Logic Expressions for a NAND Gate
- The Boolean expression for the output of a 2-input NAND gate is

X = AB where a bar over a variable or variables indicates an inversion.

13

#### The NOR Gate

- The NOR gate, like the NAND gate, is a useful logic element because it can also be used as a universal gate.
- The NOR is the same as the OR except the output is inverted.



• Operation of a NOR Gate

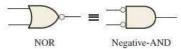
Inp	uts	Output
A	B	X
0	0	1
0	1	0
1	0	0
1	1	0





#### The NOR Gate

- Negative-AND Equivalent Operation of the NOR Gate
- For a 2-input NOR gate performing a negative-AND operation, output X is HIGH only when both inputs A and B are LOW.



• Logic Expressions for a NOR Gate

The Boolean expression for the output of a 2-input NOR gate can be written as

$$X = A + B$$

$$0 \quad 0 \quad \overline{0+0} = \overline{0} = 1$$

$$0 \quad 1 \quad 0 + \overline{1} = 0$$

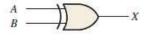
$$1 \quad 0 \quad 1 + \overline{0} = \overline{1} = 0$$

$$1 \quad 1 \quad 1 + \overline{1} = \overline{1} = 0$$

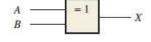
15

#### The Exclusive-OR Gate

• The output of an exclusive-OR gate is HIGH *only* when the two inputs are at opposite logic levels.



(a) Distinctive shape



(b) Rectangular outline





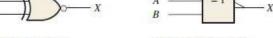
### The Exclusive-NOR Gate

• For an exclusive-NOR gate, output X is LOW when input A is LOW and input B is HIGH, or when A is HIGH and B is LOW; X is HIGH when A and B are both HIGH or both LOW.

Truth table for an exclusive-NOR gate.

Inp	outs	Output	
A	В	X	
0	0	1	
0	1	0	
1	0	0	
1	1	1	





(a) Distinctive shape

(b) Rectangular outline

17

#### **Timing Diagrams**

- a timing diagram is basically a graph that accurately displays the relationship of two or more waveforms with respect to each other on a time basis.
- Example:
- A waveform is applied to an inverter in the figure below.



- Determine the output waveform corresponding to the input and show the timing diagram.
- Solution:
- The output waveform is exactly opposite to the input (inverted).







### **Boolean Algebra and Logic Simplification**

- Boolean algebra is the mathematics of digital logic.
- A variable is a symbol (usually an italic uppercase letter or word) used to represent an action, a condition, or data. In Boolean algebra, any single variable can have only a 1 or a 0 value.
- The **complement** is the inverse of a variable and it is indicated by a bar over the variable (overbar). For example, the complement of the variable A is A. If A = 1, then A = 0. If A = 0, then A = 1.
- The **complement** of the variable *A* is read as "not *A*" or "*A* bar." Sometimes a prime symbol rather than an overbar is used to denote the complement of a variable; for example, *A'* indicates a complement of *A*.
- A literal is a variable or the complement of a variable.

19

#### **Boolean Addition**

- The **Boolean addition** is equivalent to the **OR** operation.
- In Boolean algebra, a sum term is a sum of literals.
- Some examples of sum terms are A + B,  $A + \overline{B}$ ,  $A + B + \overline{C}$ , and  $\overline{A} + B + C + D$ .
- A **sum term** is equal to 1 when one or more of the literals in the term are 1. A **sum term** is equal to 0 only if each of the literals is 0.
- Example: Determine the values of A, B, C, and D that make the sum term A + B + C + D equal to 0.
- Solution:
- For the sum term to be 0, each\_of the literals in the term must be 0. Therefore,  $A = \mathbf{0}$ ,  $B = \mathbf{1}$  so that B = 0,  $C = \mathbf{0}$ , and  $D = \mathbf{1}$  so that D = 0.
- A + B + C + D = 0 + 1 + 0 + 1 = 0 + 0 + 0 + 0 = 0





#### **Boolean Multiplication**

- Boolean multiplication is equivalent to the AND operation.
- In Boolean algebra, a product term is the product of literals.
- Some examples of product terms are  $A\overline{B}$ , AB, ABC, and  $A\overline{B}C\overline{D}$ .
- A product term is equal to 1 only if each of the literals in the term is 1. A product term is equal to 0 when one or more of the literals are 0.
- Example: Determine the values of A, B, C, and D that make the product term ABCD equal to 1.
- Solution:
- For the product term to be 1, each of the literals in the term must be 1. Therefore, A = 1, B = 0 so that B = 1, C = 1, and D = 0 so that D = 1.
  - *ABCD* = 1 . 0 . 1 . 0 = 1 . 1 . 1 . 1 = 1

21

### Laws and Rules of Boolean Algebra

• The *commutative law of addition* for two variables is written as:

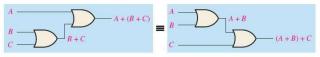
$$A+B=B+A$$

• The *commutative law of multiplication* for two variables is:

#### AB=BA

• The associative law of addition is written as follows for three variables:

$$A + (B + C) = (A+B)+C$$



The *associative law of multiplication* is written as follows for three variables:

A(BC)=(AB)C

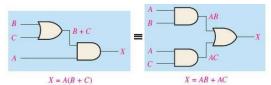




## Laws and Rules of Boolean Algebra

- Distributive Law
- The distributive law is written for three variables as follows:

#### A(B+C)=AB+AC



#### Basic rules of Boolean algebra

```
      1. A + 0 = A
      7. A \cdot A = A

      2. A + 1 = 1
      8. A \cdot \overline{A} = 0

      3. A \cdot 0 = 0
      9. \overline{\overline{A}} = A

      4. A \cdot 1 = A
      10. A + AB = A

      5. A + A = A
      11. A + \overline{AB} = A + B

      6. A + \overline{A} = 1
      12. (A + B)(A + C) = A + BC
```

A, B, or C can represent a single variable or a combination of variables.

23

## Rules of Boolean Algebra

- Rule 1: A + 0 = A. A variable ORed with 0 is always equal to the variable.
- Rule 2: A +1= 1. A variable ORed with 1 is always equal to 1.
- Rule 3: A.0= 0. A variable ANDed with 0 is always equal to 0.
- Rule 4: A. 1=A. A variable ANDed with 1 is always equal to the variable.
- Rule 5: A+A= A. A variable ORed with itself is always equal to the variable.
- Rule 6: A+A= 1. A variable ORed with its complement is always equal to 1.
- Rule 7: A.A = A. A variable ANDed with itself is always equal to the variable.
- Rule 8: A.A=0. A variable ANDed with its complement is always equal to 0.
- Rule 9: A= A. The double complement of a variable is always equal to the variable.





## Rules of Boolean Algebra

• Rule 10: A + AB = A. This rule can be proved by applying the distributive law, rule 2, and rule 4 as follows:

$$A + AB = A \cdot 1 + AB = A(1 + B)$$
 Factoring (distributive law)  
=  $A \cdot 1$  Rule 2:  $(1 + B) = 1$   
=  $A$  Rule 4:  $A \cdot 1 = A$ 

	A + AB	AB	В	A
A -	0	0	0	0
	0	0	1	0
$B \longrightarrow$	1	0	0	1
1	1	1	1	1
A straight connection	•		equ	<b>†</b>

25

## Rules of Boolean Algebra

• Rule 11:  $A + \overline{AB} = A + B$ . This rule can be proved as follows:

$$A + \overline{A}B = (A + AB) + \overline{A}B$$
 Rule 10:  $A = A + AB$   
 $= (AA + AB) + \overline{A}B$  Rule 7:  $A = AA$   
 $= AA + AB + A\overline{A} + \overline{A}B$  Rule 8: adding  $A\overline{A} = 0$   
 $= (A + \overline{A})(A + B)$  Factoring  
 $= 1 \cdot (A + B)$  Rule 6:  $A + \overline{A} = 1$   
 $= A + B$  Rule 4: drop the 1

A	В	$\overline{A}B$	$A + \overline{A}B$	A + B	_
0	0	0	0	0	$A \rightarrow \square$
0	1	1	1	1	
1	0	0	1	1	В
4	1	0	1	1	<b>↓</b>





### Rules of Boolean Algebra

• Rule 12: (A + B)(A + C) = A + BC. This rule can be proved as follows:

$$(A + B)(A + C) = AA + AC + AB + BC$$
 Distributive law  
 $= A + AC + AB + BC$  Rule 7:  $AA = A$   
 $= A(1 + C) + AB + BC$  Factoring (distributive law)  
 $= A \cdot 1 + AB + BC$  Rule 2:  $1 + C = 1$   
 $= A(1 + B) + BC$  Factoring (distributive law)  
 $= A \cdot 1 + BC$  Rule 2:  $1 + B = 1$   
 $= A + BC$  Rule 4:  $A \cdot 1 = A$ 

0 0 0 0			A + C	A + B	C	В	A
	0	0	0	0	0	0	0
0 1 0 0 0 A	0	0	1	0	1	0	0
1 0 0 0 0 -	0	0	0	1	0	1	0
1 1 1 1 C	1	1	1	1	1	1	0
1 1 1 0 1	0	1	1	1	0	0	1
1 1 1 0 1	0	1	1	1	1	0	1
1 1 1 0 1 1	0	1	1	1	0	1	1
1 1 1 1 1 1 0	1	1	1	1	1	1	1
equal	Taxable 1	Ť					

27

### DeMorgan's Theorems

- DeMorgan's first theorem is stated as follows:
- The complement of a product of variables is equal to the sum of the complements of the variables.

The formula for expressing this theorem for two variables is

$$\overline{XY} = \overline{X} + \overline{Y}$$

DeMorgan's second theorem is stated as follows:

• The complement of a sum of variables is equal to the product of the complements of the variables.

The formula for expressing this theorem for two variables is

$$\overline{X + Y} = \overline{X}\overline{Y}$$





#### Applying DeMorgan's Theorems

• The following procedure illustrates the application of DeMorgan's theorems and Boolean algebra to the specific expression

$$\overline{A + B\overline{C}} + D(\overline{E + F})$$

Step 1: Identify the terms to which you can apply DeMorgan's theorems, and think of each term as a single variable. Let  $\overline{A + B\overline{C}} = X$  and  $D(\overline{E + \overline{F}}) = Y$ .

Step 2: Since  $\overline{X} + \overline{Y} = \overline{X}\overline{Y}$ ,

$$\overline{(\overline{A} + B\overline{\overline{C}}) + (\overline{D(E + \overline{F})})} = (\overline{\overline{A} + B\overline{\overline{C}}})(\overline{D(E + \overline{F})})$$

**Step 3:** Use rule  $9(\overline{\overline{A}} = A)$  to cancel the double bars over the left term (this is not part of DeMorgan's theorem).

$$(\overline{\overline{A} + B\overline{\overline{C}}})(\overline{D(\overline{E} + \overline{F})}) = (A + B\overline{C})(\overline{D(\overline{E} + \overline{F})})$$

Step 4: Apply DeMorgan's theorem to the second term.

$$(A+B\overline{C})(\overline{D(\overline{E}+\overline{F})})=(A+B\overline{C})(\overline{D}+(\overline{\overline{E}+\overline{F}}))$$

**Step 5:** Use rule  $9(\overline{\overline{A}} = A)$  to cancel the double bars over the  $E + \overline{F}$  part of the term.

$$(A + B\overline{C})(\overline{D} + \overline{E + F}) = (A + B\overline{C})(\overline{D} + E + \overline{F})$$

29

## Applying DeMorgan's Theorems

• Example: Apply DeMorgan's theorems to the following expression:

$$\overline{(A+B+C)D}$$

• Solution:

Let A + B + C = X and D = Y. The expression  $\overline{(A + B + C)D}$  is of the form  $\overline{XY} = \overline{X} + \overline{Y}$  and can be rewritten as

$$\overline{(A+B+C)D} = \overline{A+B+C} + \overline{D}$$

Next, apply DeMorgan's theorem to the term  $\overline{A + B + C}$ .

$$\overline{A+B+C}+\overline{D}=\overline{A}\overline{B}\overline{C}+\overline{D}$$





## Applying DeMorgan's Theorems

• Example: Apply DeMorgan's theorems to the following expression:

$$\overline{ABC + DEF}$$

#### Solution

Let ABC = X and DEF = Y. The expression  $\overline{ABC + DEF}$  is of the form  $\overline{X + Y} = \overline{XY}$  and can be rewritten as

$$\overline{ABC} + \overline{DEF} = (\overline{ABC})(\overline{DEF})$$

Next, apply DeMorgan's theorem to each of the terms  $\overline{ABC}$  and  $\overline{DEF}$ .

$$(\overline{ABC})(\overline{DEF}) = (\overline{A} + \overline{B} + \overline{C})(\overline{D} + \overline{E} + \overline{F})$$

31

## Applying DeMorgan's Theorems

• Example: Apply DeMorgan's theorems to the following expression:

$$\overline{AB} + \overline{CD} + EF$$

Solution:

Let  $A\overline{B} = X$ ,  $\overline{C}D = Y$ , and EF = Z. The expression  $\overline{AB} + \overline{C}D + EF$  is of the form  $\overline{X} + \overline{Y} + \overline{Z} = \overline{X}\overline{Y}\overline{Z}$  and can be rewritten as

$$\overline{A\overline{B}} + \overline{C}D + EF = (\overline{A\overline{B}})(\overline{\overline{C}D})(\overline{EF})$$

Next, apply DeMorgan's theorem to each of the terms  $\overline{AB}$ ,  $\overline{\overline{CD}}$ , and  $\overline{EF}$ .

$$(\overline{A}\overline{B})(\overline{\overline{C}D})(\overline{EF}) = (\overline{A} + B)(C + \overline{D})(\overline{E} + \overline{F})$$





#### Applying DeMorgan's Theorems

- Example: Apply DeMorgan's theorems to each of the following expressions:
  - (a)  $\overline{(\overline{A+B})} + \overline{\overline{C}}$
  - (b)  $\overline{(\overline{A} + B) + CD}$
  - (c)  $\overline{(A+B)\overline{C}D} + E + \overline{F}$
- Solution
- (a)  $\overline{(\overline{A}+\overline{B})}+\overline{\overline{C}}=(\overline{\overline{A}+\overline{B}})\overline{\overline{\overline{C}}}=(A+B)C$
- (b)  $\overline{(\overline{A}+B)+CD}=(\overline{\overline{A}+B})\overline{CD}=(\overline{\overline{A}B})(\overline{C}+\overline{D})=A\overline{B}(\overline{C}+\overline{D})$
- (c)  $\overline{(A+B)\overline{C}D} + E + \overline{F} = \overline{((A+B)\overline{C}D)}(\overline{E} + \overline{F}) = (\overline{A}\overline{B} + C + D)\overline{E}F$

33

## Applying DeMorgan's Theorem

- Example:
- The Boolean expression for an exclusive-OR gate is AB + AB. With this as a starting point, use DeMorgan's theorems and any other rules or laws that are applicable to develop an expression for the exclusive-NOR gate.
- Solution:

Start by complementing the exclusive-OR expression and then applying DeMorgan's theorems as follows:

$$\overline{A\overline{B} + \overline{A}B} = (\overline{A\overline{B}})(\overline{\overline{A}B}) = (\overline{A} + \overline{\overline{B}})(\overline{\overline{A}} + \overline{B}) = (\overline{A} + B)(A + \overline{B})$$

Next, apply the distributive law and rule 8 ( $A \cdot \overline{A} = 0$ ).

$$(\overline{A} + B)(A + \overline{B}) = \overline{A}A + \overline{A}\overline{B} + AB + B\overline{B} = \overline{A}\overline{B} + AB$$

The final expression for the XNOR is  $\overline{AB} + AB$ . Note that this expression equals 1 any time both variables are 0s or both variables are 1s.