

جام<u>عة</u> الم<u>ستقبل</u> AL MUSTAQBAL UNIVERSITY

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Department of Cyber Security

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Class: Second

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Lecture: (3)

Queue - 2

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CIRCULAR QUEUE

A more efficient queue representation is obtained by regarding the array Q[MAX] as circular. Any number of items could be placed on the queue. This implementation of a queue is called a circular queue because it uses its storage array as if it were a circle instead of a linear list.

There are two problems associated with linear queue. They are:

- Time consuming: linear time to be spent in shifting the elements to the beginning of the queue.
- Signaling queue full: even if the queue is having vacant position.

For example, let us consider a linear queue status as follows:



Next insert another element, say 66 to the queue. We cannot insert 66 to the queue as the rear crossed the maximum size of the queue (i.e., 5). There will be queue full signal. The queue status is as follows:



This difficulty can be overcome if we treat queue position with index zero as a position that comes after position with index four then we treat the queue as a **circular queue**.

In circular queue if we reach the end for inserting elements to it, it is possible to insert new elements if the slots at the beginning of the circular queue are empty.

Representation of Circular Queue:

Let us consider a circular queue, which can hold maximum (MAX) of six elements. Initially the queue is empty.



Now, insert 11 to the circular queue. Then circular queue status will be:



Insert new elements 22, 33, 44 and 55 into the circular queue. The circular queue status is:



Now, delete an element. The element deleted is the element at the front of the circular queue. So, 11 is deleted. The circular queue status is as follows:



Again, delete an element. The element to be deleted is always pointed to by the FRONT pointer. So, 22 is deleted. The circular queue status is as follows:



Again, insert another element 66 to the circular queue. The status of the circular queue is:



Now, insert new elements 77 and 88 into the circular queue. The circular queue status is:



Now, if we insert an element to the circular queue, as COUNT = MAX we cannot add the element to circular queue. So, the circular queue is *full*.

Operations on Circular queue:

a.enqueue() or insertion():This function is used to insert an element into the circular queue. In a circular queue, the new element is always inserted at Rear position.

void insertCQ()	Algorithm: procedure of insertCQ():		
{			
int data;	Step-1:START		
if(count ==MAX)	Step-2: if count==MAX then		
{	Write "Circular queue is full"		
printf("\n Circular Queue is Full");	Step-3:otherwise		
}	3.1: read the data element		
else	3.2: CQ[rear]=data		
{	3.3: rear=(rear+1)%MAX		
printf("\n Enter data: ");	3.4: count=count+1		
scanf("%d", &data);	Step-4:STOP		
CQ[rear] = data;			
rear = (rear + 1) % MAX;			
count ++;			
printf("\n Data Inserted in the Circular			
Queue ");			
}			
}			

b.dequeue() or deletion():This function is used to delete an element from the circular queue. In a circular queue, the element is always deleted from front position.

void deleteCQ()	Algorithm: procedure of deleteCQ():
{	
if(count ==0)	Step-1:START
{	Step-2: if count==0 then
<pre>printf("\n\nCircular Queue is Empty");</pre>	Write "Circular queue is empty"
}	Step-3:otherwise
else	3.1: print the deleted element
{	3.2: front=(front+1)%MAX
printf("\n Deleted element from Circular	3.3: count=count-1
Queue is %d ", CQ[front]);	Step-4:STOP
front = (front + 1) % MAX;	
count;	
}	
}	

c.dispaly():This function is used to display the list of elements in the circular queue.

void displayCQ()	Algorithm: procedure of displayCQ():
{	
int i, j;	Step-1:START
if(count ==0)	Step-2: if count==0 then
{	Write "Circular queue is empty"
<pre>printf("\n\n\t Circular Queue is Empty ");</pre>	Step-3:otherwise
}	3.1: print the list of elements
else	3.2: for i=front to j!=0
{	3.3: print CQ[i]
printf("\n Elements in Circular Queue are:	3.4: i=(i+1)%MAX
");	Step-4:STOP
j = count;	
for(i = front; j != 0; j)	
{	
printf("%d\t", CQ[i]);	
i = (i + 1) % MAX;	
}	
}	
}	

Deque:

In the preceding section we saw that a queue in which we insert items at one end and from which we remove items at the other end. In this section we examine an extension of the queue, which provides a means to insert and remove items at both ends of the queue. This data structure is a *deque*. The word *deque* is an acronym derived from *double-ended queue*. Below figure shows the representation of a deque.



Figure Representation of a deque.

deque provides four operations. Below Figure shows the basic operations on a deque.

- enqueue_front: insert an element at front.
- dequeue_front: delete an element at front.
- enqueue_rear: insert element at rear.
- dequeue_rear: delete element at rear.



Figure * * Basic operations on deque

There are two variations of deque. They are:

- Input restricted deque (IRD)
- Output restricted deque (ORD)

An Input restricted deque is a deque, which allows insertions at one end but allows deletions at both ends of the list.

An output restricted deque is a deque, which allows deletions at one end but allows insertions at both ends of the list.

Priority Queue:

A **priority queue** is a collection of elements such that each element has been assigned a priority. We can insert an element in priority queue at the rare position. We can delete an element from the priority queue based on the elements priority and such that the order in which elements are deleted and processed comes from the following rules:

1. An element of higher priority is processed before any element of lower priority.

2. Two elements with same priority are processed according to the order in which they were added to the queue. It follows FIFO or FCFS(First Comes First serve) rules.

We always remove an element with the highest priority, which is given by the minimal integer priority assigned.

[3] [1] [4] [2] [₽] [5]					priority	
	5	10	30	25	40	Queue
[0] [1] [2] [3] [4]					index	

A prototype of a priority queue is time sharing system: programs of high priority are processed first, and programs with the same priority form a standard queue. An efficient implementation for the Priority Queue is to use heap, which in turn can be used for sorting purpose called heap sort

Priority queues are two types:

- 1. Ascending order priority queue
- 2. Descending order priority queue

1. Ascending order priority queue: It is Lower priority number to high priority number. Examples: order is 1,2,3,4,5,6,7,8,9,10

2. Descending order priority queue: It is high priority number to lowest priority number. Examples: Order is 10,9,8,7,6,5,4,3,2,1

Implementation of Priority Queue:

Implementation of priority queues are two types:

- 1. Through Queue(Using Array)
- 2. Through Sorted List(Using Linked List)

1. Through Queue (Using Array): In this case element is simply added at the rear end as usual. For deletion, the element with highest priority is searched and then deleted.



2. Through sorted List (Using Linked List): In this case insertion is costly because the element insert at the proper place in the list based on the priority. Here deletion is easy since the element with highest priority will always be in the beginning of the list.

Insert 40 priority 1



1. Difference between stacks and Queues?				
stacks	Queues			
1.A stack is a linear list of elements in which	1.A Queue is a linerar list of elements in which			
the element may be inserted or deleted at	the elements are added at one end and			
one end.	deletes the elements at another end.			
	2 In Queue the element which is inserted			
2. In stacks, elements which are inserted last is the first element to be deleted.	first is the element deleted first.			
	3. Queues are called FIFO (First In First			
3.Stacks are called LIFO (Last In First Out)list	Out)list.			
	4. In Queue elements are removed in the			
4.In stack elements are removed in reverse order in which thy are inserted.	same order in which thy are inserted.			
	5. Suppose the elements a,b,c,d,e are inserted			
5.suppose the elements a,b,c,d,e are	in the Queue, the deletion of elements will be			
inserted in the stack, the deletion of elements will be e,d,c,b,a.	in the same order in which thy are inserted.			
	6. In Queue there are two pointers one for			
6.In stack there is only one pointer to insert and delete called "Top".	insertion called "Rear" and another for deletion called "Front".			
7.Initially top=-1 indicates a stack is empty.	7. Initially Rear=Front=-1 indicates a Queue is empty.			
8.Stack is full represented by the condition				
TOP=MAX-1(if array index starts from '0').	8.Queue is full represented by the condition Rear=Max-1.			
9.To push an element into a stack, Top is				
incremented by one	9.To insert an element into Queue, Rear is incremented by one.			
10.To POP an element from stack,top is				
decremented by one.	10.To delete an element from Queue. Front is			

