



جامعة المستقبل
AL MUSTAQBAL UNIVERSITY



قسم الامن السيبراني
DEPARTMENT OF CYBER SECURITY

SUBJECT:

SEARCHING AND SORTING ALGORITHMS

CLASS:

SECOND

LECTURER:

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LECTURE: (5-2)

POST-ORDER TRAVERSAL



Program to Implement Postorder Traversal of Binary Tree

Below is the code implementation of the postorder traversal:

```
// C++ program for postorder traversals

#include <bits/stdc++.h>
using namespace std;

// Structure of a Binary Tree Node
struct Node {
    int data;
    struct Node *left, *right;
    Node(int v)
    {
        data = v;
        left = right = nullptr;
    }
};

// Function to print postorder traversal
void printPostorder(struct Node* node)
{
    if (node == nullptr)
        return;

    // First recur on left subtree
    printPostorder(node->left);

    // Then recur on right subtree
    printPostorder(node->right);

    // Now deal with the node
    cout << node->data << " ";
}

// Driver code
int main()
{
    struct Node* root = new Node(1);
    root->left = new Node(2);
    root->right = new Node(3);
    root->left->left = new Node(4);
    root->left->right = new Node(5);
}
```



```
root->right->right = new Node(6);

// Function call
cout << "Postorder traversal of binary tree is: \n";
printPostorder(root);

return 0;
}
```

Output

Postorder traversal of binary tree is:

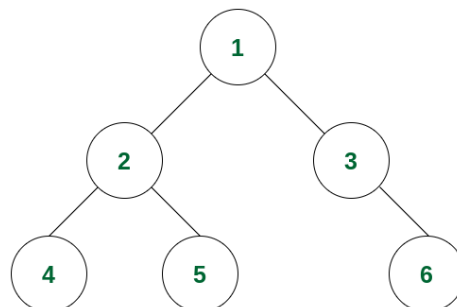
4 5 2 6 3 1

- **Complexity Analysis:**

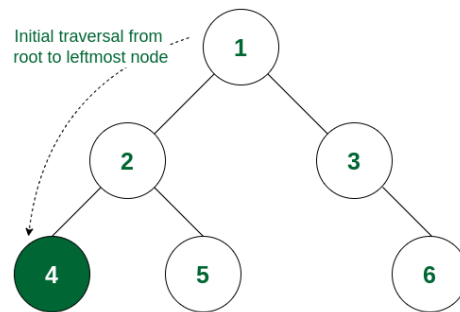
- ❖ **Time Complexity:** $O(N)$ where N is the total number of nodes. Because it traverses all the nodes at least once.
- ❖ **Auxiliary Space:** $O(1)$ if no recursion stack space is considered. Otherwise, $O(h)$ where h is the height of the tree
 - In the worst case, h can be the same as N (when the tree is a skewed tree)
 - In the best case, h can be the same as $\log N$ (when the tree is a complete tree)

How does Postorder Traversal of Binary Tree Work?

Consider the following tree:

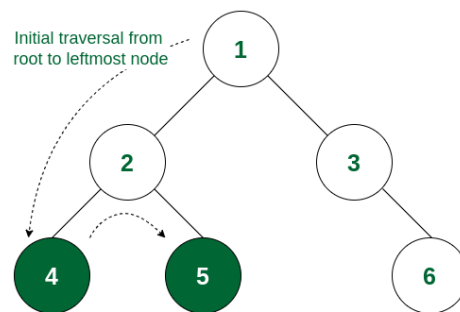


- ✓ **Step 1:** The traversal will go from 1 to its left subtree i.e., 2, then from 2 to its left subtree root, i.e., 4. Now 4 has no subtree, so it will be visited.



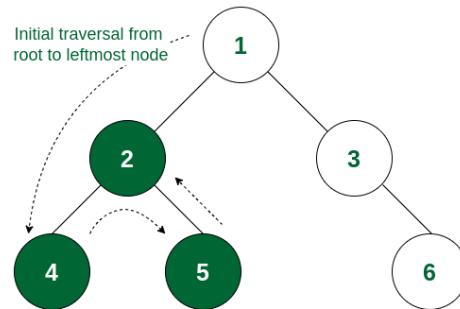
The leftmost leaf node (i.e., 4) is visited first

- ✓ **Step 2:** As the left subtree of 2 is visited completely, now it will traverse the right subtree of 2 i.e., it will move to 5. As there is no subtree of 5, it will be visited.



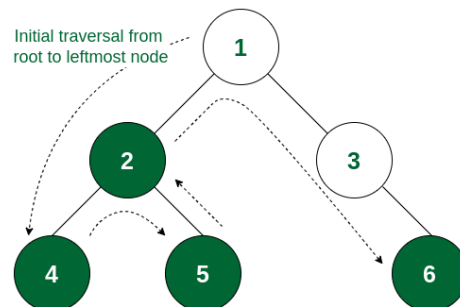
Left subtree of 2 is traversed. So 5 is visited next

- ✓ **Step 3:** Now both the left and right subtrees of node 2 are visited. So now visit node 2 itself.



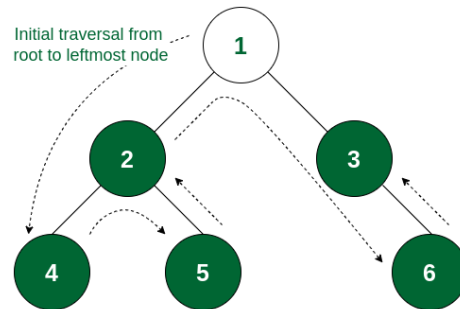
All subtrees of 2 are visited. So 2 is visited next

- ✓ **Step 4:** As the left subtree of node 1 is traversed, it will now move to the right subtree root, i.e., 3. Node 3 does not have any left subtree, so it will traverse the right subtree i.e., 6. Node 6 has no subtree and so it is visited.



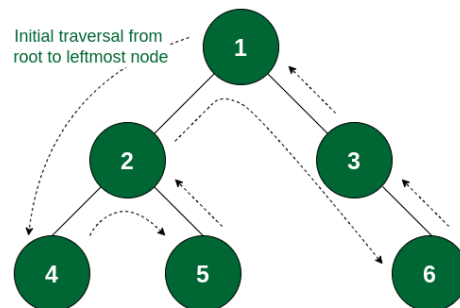
6 has no subtrees. So it is visited

- ✓ **Step 5:** All the subtrees of node 3 are traversed. So now node 3 is visited.



3 is visited after all its subtrees are traversed

- ✓ **Step 6:** As all the subtrees of node 1 are traversed, now it is time for node 1 to be visited and the traversal ends after that as the whole tree is traversed.



The root of the tree (i.e., 1) is visited

- ✓ So the order of traversal of nodes is 4 -> 5 -> 2 -> 6 -> 3 -> 1.