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Lecture: (5)

Blind Search

Subject: Artificial Intelligence Principles
Level: Third
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1. Problem Solving using Search

The breadth-first algorithm spreads out in a uniform manner from the start node. From the start, it looks at each node one edge away. Then it moves out from those nodes to all nodes two edges away from the start. This continues until either the goal node is found or the entire tree is searched.

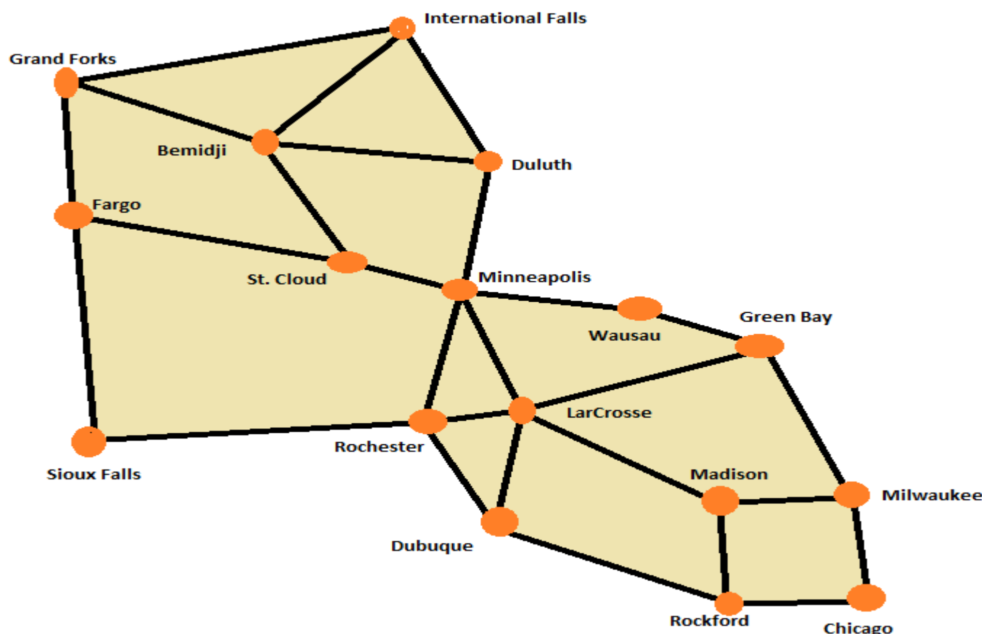
2. Characteristics of breadth-first algorithm

- Breadth-first search is ***complete***; It will find a solution if one exists.
- But it is neither ***optimal*** in the general case (it won't find the best solution, just the first one that matches the goal state),
- It doesn't have ***good time or space complexity*** (it grows exponentially in time and memory consumption).

Example

A map like the one in Figure below can be naturally represented by *a graph data structure*, where the cities names are the nodes, and the major roadways between cities are the *links or edges* of the graph. So, from a programming perspective, our problem is to traverse a graph data structure in a systematic way until we either find the goal city or exhaust all possibilities. Hopefully having the entire state-space shown on a map will make understanding the operations of the search algorithms easier. In more complex problems, *all we have is the single start state and a set of operators which are used to generate more and more new states*. The search algorithms work the same way, but conceptually, *we are growing or expanding the graph, instead of having it specified at the start.* Map of midwestern U.S. cities is illustrated below:-

The *breadth-first algorithm* spreads out in a uniform manner from the start node. From the start, it looks at each node one edge away. Then it moves out from those nodes to all nodes two edges away from the start. This continues until either the goal node is found or the entire tree is searched.



Let's walk through an example to see how breadth-first search could find a city on our map.

1- Our search begins in **Rochester (Start State)**, and we want to know if we can get to **Wausau (Goal State)** from there.

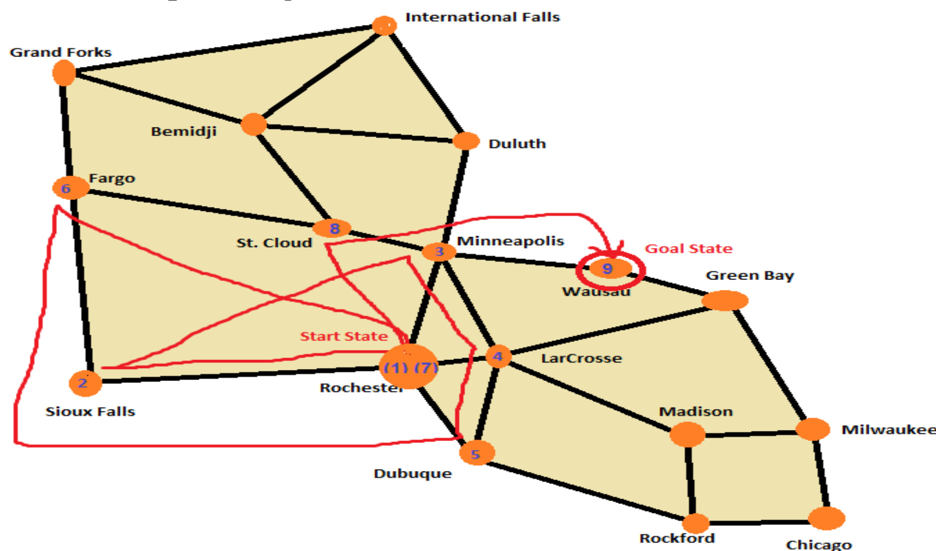
2- The **Rochester** node is placed on the *queue* in step 1 in previous algorithm. Next we enter our search loop at step 2. Queue=[**Rochester]**

3- We remove Rochester, the first node from the queue. **Rochester** does not contain our goal state (**Wausau**) so we expand it by taking each child node in **Rochester**, *and adding them to the back of the queue.* Queue= **[Sioux Falls, Minneapolis, LaCrosse, and Dubuque].**

4- We remove the first node from the queue (**Sioux Falls**) and test it to see if it is our goal state. It is not, so we expand it, adding ***Fargo*** and ***Rochester*** to the end of our queue, which now contains **[Minneapolis, LaCrosse, Dubuque, Fargo, and Rochester]**.

5- We remove **Minneapolis**, the goal test fails, and we expand that node, adding **St.Cloud, Wausau, Duluth, LaCrosse**, and ***Rochester*** to the search queue, now holding **[LaCrosse, Dubuque, Fargo, Rochester, St.Cloud, Wausau, Duluth, LaCrosse, and Rochester]**.
6- We test **LaCrosse** and then expand it, adding **Minneapolis, GreenBay, Madison, Dubuque**, and **Rochester** to the list, which has now grown to **[Dubuque, Fargo, Rochester, St.Cloud, Wausau, Duluth, LaCrosse, Rochester, Minneapolis, GreenBay, Madison, Dubuque, and Rochester]**. We remove **Dubuque** and add **Rochester, LaCrosse**, and **Rockford** to the search queue.
7- At this point, we have tested every node which is one level in the tree away from the start node (**Rochester**). Our search queue contains the following nodes: **[Fargo, Rochester, St.Cloud, Wausau, Duluth, LaCrosse, Rochester, Minneapolis, GreenBay, Madison, Dubuque, Rochester, Rochester, LaCrosse, and Rockford]**.
8- We remove **Fargo**, which is two levels away from **Rochester**, and add **Grand Forks, St. Cloud**, and **Sioux Falls**.
9- Then we test and expand **Rochester** (Rochester to Minneapolis to Rochester is two levels away from our start). Next is **St. Cloud**; again we expand that node.

10- Finally, we get to **Wausau**; our goal test succeeds and we declare success.
11- Our search order was Rochester, Sioux Falls, Minneapolis, LaCrosse, Dubuque, Fargo, Rochester, St. Cloud, and Wausau as shown below:-



Note that this trace could have been greatly simplified by ***keeping track of nodes which had been tested and expanded***. This would have cut down on our time and space complexity.

3. Characteristics of Depth First Search

- The depth-first algorithm searches from the start or root node all the way down to a leaf node. If it does not find the goal node, ***it backtracks up the tree and searches down the next untested path until it reaches the next leaf***.

- If you imagine a large tree, the depth-first algorithm may spend a large amount of time searching the paths on the lower left when the answer ***is really in the lower right***.
- Depth-first search is a brute-force method, it will blindly follow this search pattern until it comes across a node containing the goal state, or it searches the entire tree.
- Depth-first search has *lower* memory requirements than breadth-first search,
- It is neither complete nor optimal.

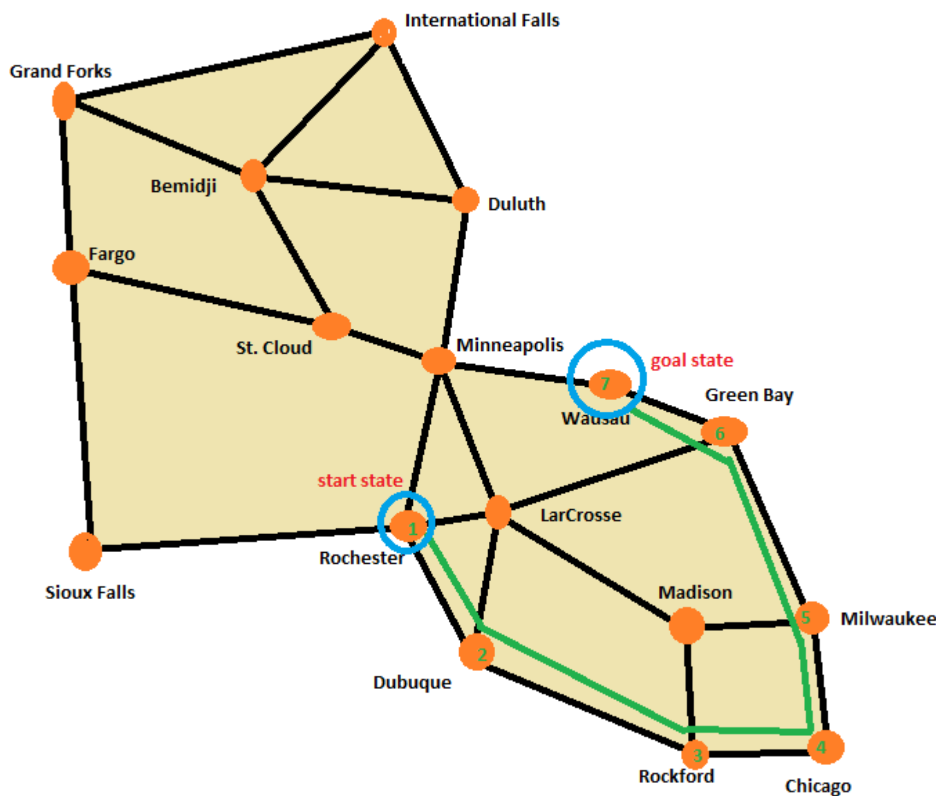
Example

Let's walk through a simple example of how depth-first search would work if we started in **Rochester** and wanted to see if we could get to **Wausau**.

Initial State= (Rochester) and Final-State (Wausau)

1. Starting with **Rochester**, we test and expand it, placing **Sioux Falls**, then **Minneapolis**, then **LaCrosse**, then **Dubuque** at the front of the search queue ***Queue***=**[Dubuque, LaCrosse, Minneapolis, Sioux Falls]**.
2. We remove **Dubuque** and test it; it fails, so we expand it adding **Rochester** to the front, then **LaCrosse**, then **Rockford**. Our search queue now looks like **[Rockford, LaCrosse, Rochester, LaCrosse, Minneapolis, Sioux Falls**].
3. We remove **Rockford**, and add **Dubuque, Madison**, and **Chicago** to the front of the queue in that order, yielding ***Queue***=**[Chicago, Madison, Dubuque, LaCrosse, Rochester, LaCrosse, Minneapolis, Sioux Falls]**.
4. We test **Chicago**, and place **Rockford**, and **Milwaukee** on the queue.
5. We take **Milwaukee** from the front and add **Chicago, Madison**, and **Green Bay** to the search queue. It is now ***Queue***=**[Green Bay, Madison, Chicago, Rockford, Chicago, Madison, Dubuque, LaCrosse, Rochester, LaCrosse, Minneapolis, Sioux Falls]**.

6. We remove **Green Bay** and add **Milwaukee, LaCrosse**, and **Wausau** to the queue in that order. Finally, **Wausau** is at the front of the queue and our goal test succeeds and our search ends. Our search order was Rochester, Dubuque, Rockford, Chicago, Milwaukee, Green Bay, and Wausau.



In this example, *we again did not prevent tested nodes from being added to the search queue*. As a result, we had duplicate nodes on the queue. In the depth-first case, this could have been disastrous. *We could have easily had a cycle or loop where we tested one city, then a second, then the first again, ad infinitum*.

4.Tree Data Structure

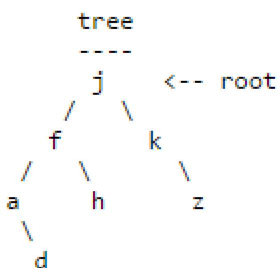
Certain programming problems are easier to solve using multiple data structures. For example, testing a sequence of characters to determine if it is a palindrome (i.e., reads the same forward and backward, like "radar") can be accomplished easily with one *stack* and one *queue*. The solution is to enter the sequence of characters into both data structures, then remove letters from each data structure one at a time and compare them, making sure that the letters match.

In this palindrome example, the user (person writing the main program) has access to both data structures to solve the problem. Another way that 2 data structures can be used in concert is to use one data structure to help implement another.

We will examine how a common data structure can be used to help traverse a tree in breadth-first order.

5. Depth-first traversal:

We have already seen a few ways to traverse the elements of a tree. For example, given the following tree:



A *preorder traversal* would visit the elements in the order: **j, f, a, d, h, k, z**.

This type of traversal is called a *depth-first* traversal. Why? Because it tries to go deeper in the tree before exploring siblings. For example, the traversal visits all the descendants of **f** (i.e., keeps going deeper) before visiting **f**'s sibling **k** (and any of **k**'s descendants).

As we've seen, this kind of traversal can be achieved by a simple recursive algorithm:

PREORDER-TRAVERSE(tree)

```
if (tree not empty)
  visit root of tree
  PREORDER-TRAVERSE(left subtree)
  PREORDER-TRAVERSE(right subtree)
```

The 2 other traversal orders we know are *inorder* and *postorder*. An inorder traversal would give us: **a, d, f, h, j, k, z**. A postorder traversal would give us: **d, a, h, f, z, k, j**.

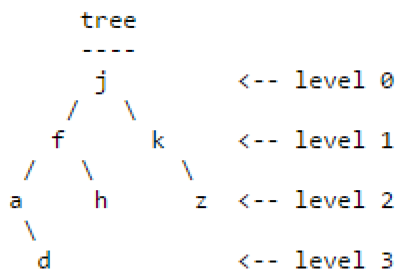
Well, inorder and postorder traversals, like a preorder traversal, also try to go *deeper* first...

For example, the inorder traversal visits **a** and **d** before it explores **a**'s sibling **h**. Likewise, it visits all of **j**'s left subtree (i.e., "a, d, f, h") before exploring **j**'s right subtree (i.e., "k, z"). The same is true for the postorder traversal. It visits all of **j**'s left subtree (i.e., "d, a, h, f") before exploring any part of the right subtree (i.e., "z, k").

6. Breadth-first traversal:

Depth-first is not the only way to go through the elements of a tree. Another way is to go through them *level-by-level*.

For example, each element exists at a certain *level* (or depth) in the tree:



(Computer people like to number things starting with 0.) So, if we want to visit the elements level-by-level (and left-to-right, as usual), we would start at level 0 with **j**, then go to level 1 for **f** and **k**, then go to level 2 for **a, h** and **z**, and finally go to level 3 for **d**.

This level-by-level traversal is called a *breadth-first traversal* because we explore the *breadth*, i.e., full width of the tree at a given level, before going *deeper*.

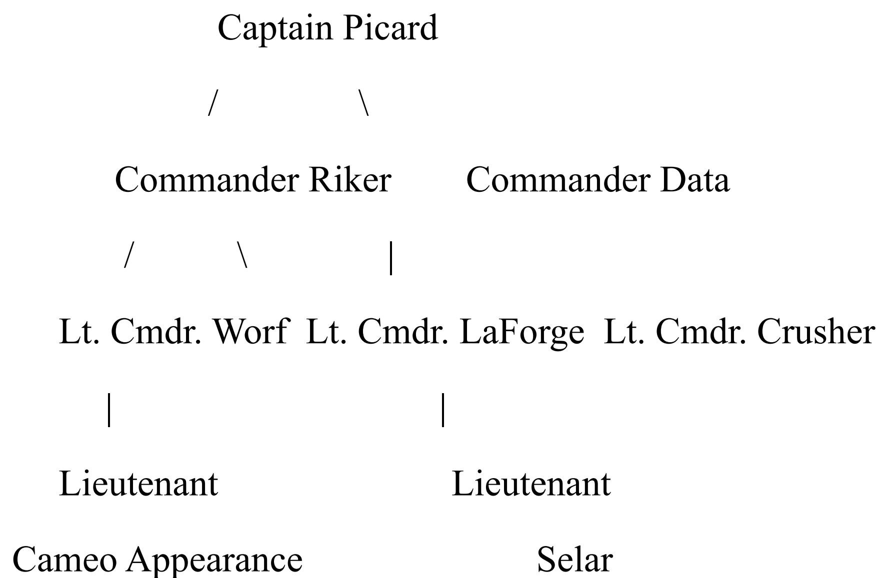
Now, how might we traverse a tree breadth-first? We'll need some other mechanism than the ones we've already used since *preorder, inorder* and *postorder* traversals don't produce breadth-first order.

7. Why breadth-first:

You may be thinking: "Why would we ever want to traverse a tree breadth-first?" Well, there are many reasons....

Tree of Officers

Suppose you have a tree representing some command structure:



This tree is meant to represent who is in charge of lower-ranking officers. For example, Commander Riker is directly responsible for Worf and LaForge. People of the same rank are at the same level in the tree. However, to distinguish between people of the same rank, those with more experience are on the left and those with less on the right (i.e., experience decreases from left to right).

Suppose a fierce battle with an enemy ensues. If officers start dropping like flies, we need to know who is the next person to take over command. One way to trace the path that command will follow is to list the officers in the tree in *breadth-first* order. This would give:

1. Captain Picard
2. Commander Riker
3. Commander Data
4. Lt. Cmdr. Worf
5. Lt. Cmdr. LaForge
6. Lt. Cmdr. Crusher
7. Lieutenant Cameo-Appearance
8. Lieutenant Selar

Let's return to example trees that are binary and that just hold characters.

As we've seen, the *recursive* tree traversals go deeper in the tree first. Instead, if we are going to implement a breadth-first traversal of a tree, we'll need some help....Perhaps one of the data structures we already know can be of assistance?

How to use helper data structure: What we'll do is store each element in the tree in a data structure and process (or *visit*) them as we *remove* them from the data structure.

We can best determine what data structure we need by looking at an example:

```
  f
 / \
a   h
\
 d
```

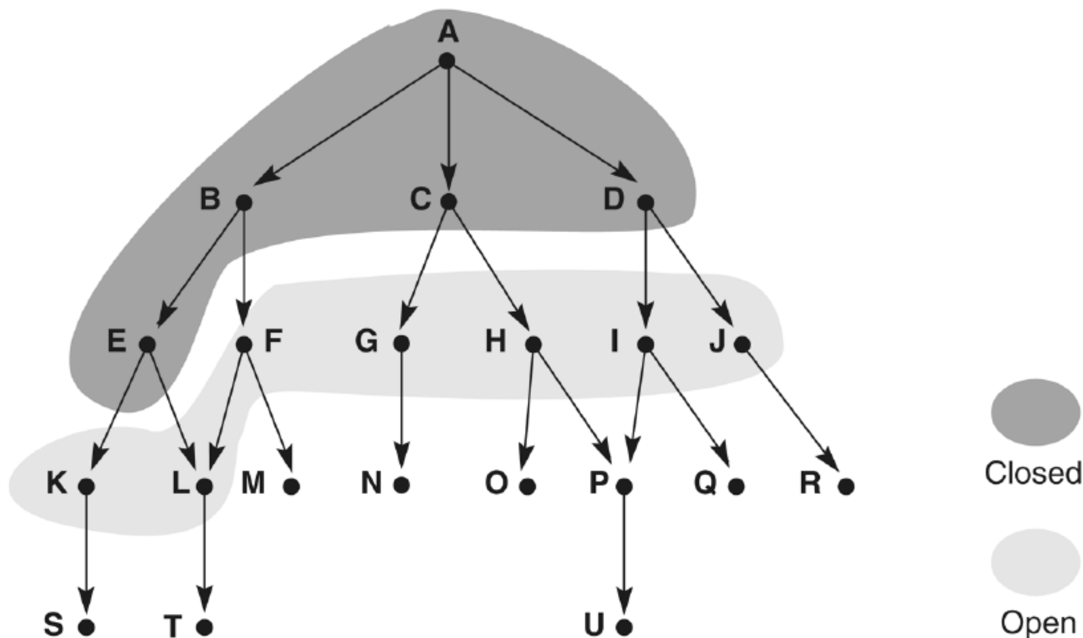
When we are at element **f**, that is the only time we have access to its 2 immediate children, **a** and **h**. So, when we are at **f**, we'd better put its children in the data

structure. Obviously then, **f** must have been in the data structure before them (i.e., first), since we'd have put **f** in when we were at **f**'s parent.

So, if we put the parent in the data structure before its children, what data structure will give us the order we need? In other words, to explore the tree breadth-first, do we want the children to be removed from the data structure first or the parent to be removed first?

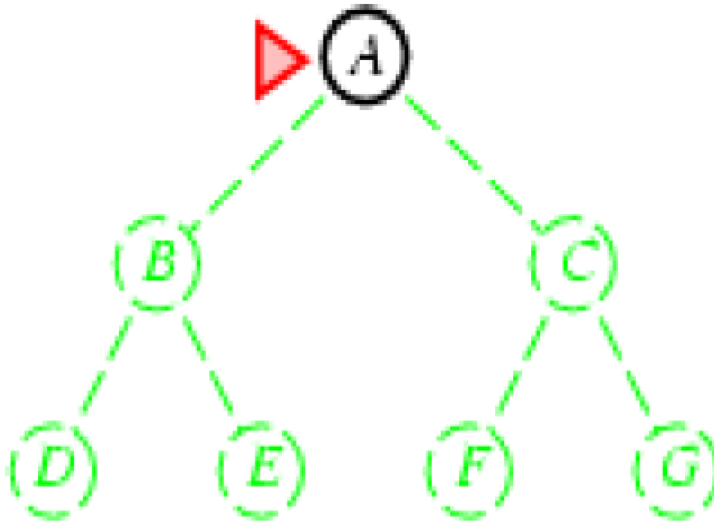
Answer: A **queue** will give us the order we want! A queue enforces *first-in-first-out* order, and we want to process the first thing in the data structure, the parent, before its descendents.

<u>Example of Breadth First Search</u>



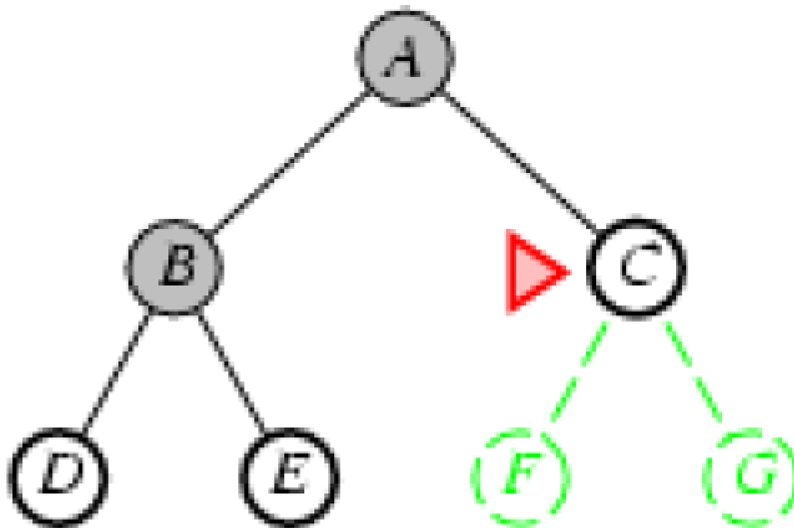
Open: [A]

Closed[ ]



Open: [B,C]

Closed:[A]

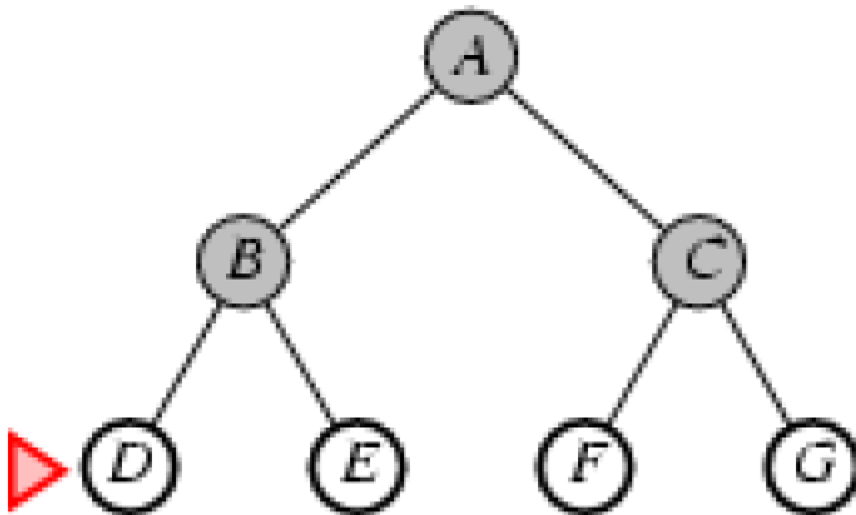


Open:[C,D,E]

Closed:[A,B]

We continue in the same procedure until we reach the Goal [G]

Time Complexity



Example 2: BFS (Breadth First Search)

Taken from http://iis.kaist.ac.kr/es/

2	8	3
1	6	4
7		5

1	2	3
8		4
7	6	5

Initial State

Goal State:

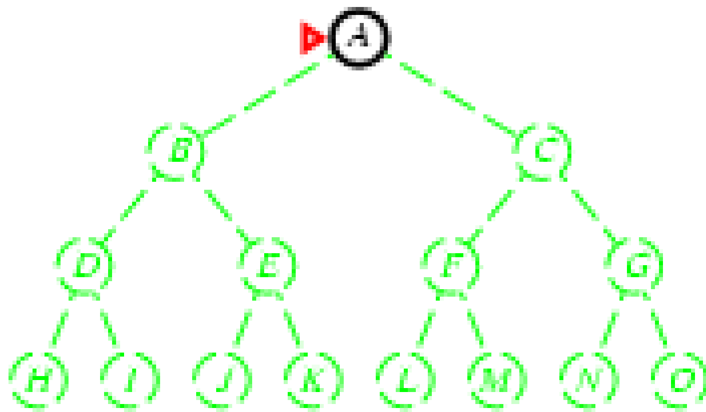
How can we solve this problem?

1. **We need to initiate the start (initial State) .**
2. **We need to use the control strategy.**
3. **Apply the initial state and control strategy to reach the Goal.**

The solving of this problem is to consider the initial state as above:

Goal

Example3: Apply this example using Depth First Search with goal [N]?



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